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Published in:
Chemical Engineering Research and Design

Link to article, DOI:
[10.1016/j.cherd.2020.03.030](https://doi.org/10.1016/j.cherd.2020.03.030)

Publication date:
2020

Document Version
Peer reviewed version

[Link back to DTU Orbit](#)

Citation (APA):
Bähner, F. D., Prado-Rubio, O. A., & Huusom, J. K. (2020). Discrete-Continuous Dynamic Simulation of Plantwide Batch Process Systems in MATLAB. *Chemical Engineering Research and Design*, 159, 66-77. <https://doi.org/10.1016/j.cherd.2020.03.030>

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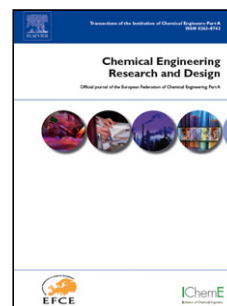
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PII: S0263-8762(20)30132-5

DOI: <https://doi.org/doi:10.1016/j.cherd.2020.03.030>

Reference: CHERD 4050

To appear in: *Chemical Engineering Research and Design*

Received Date: 16 January 2020

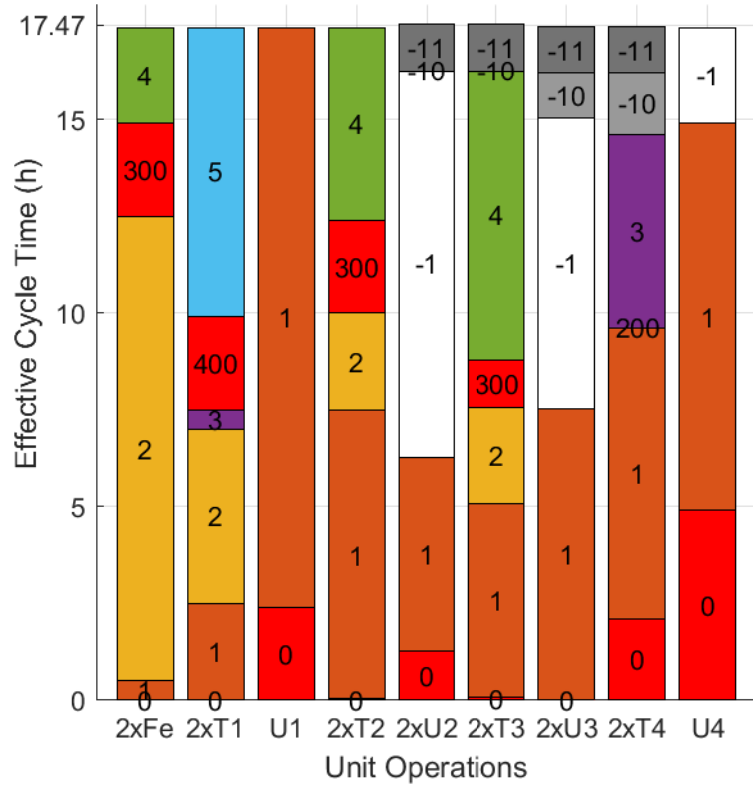
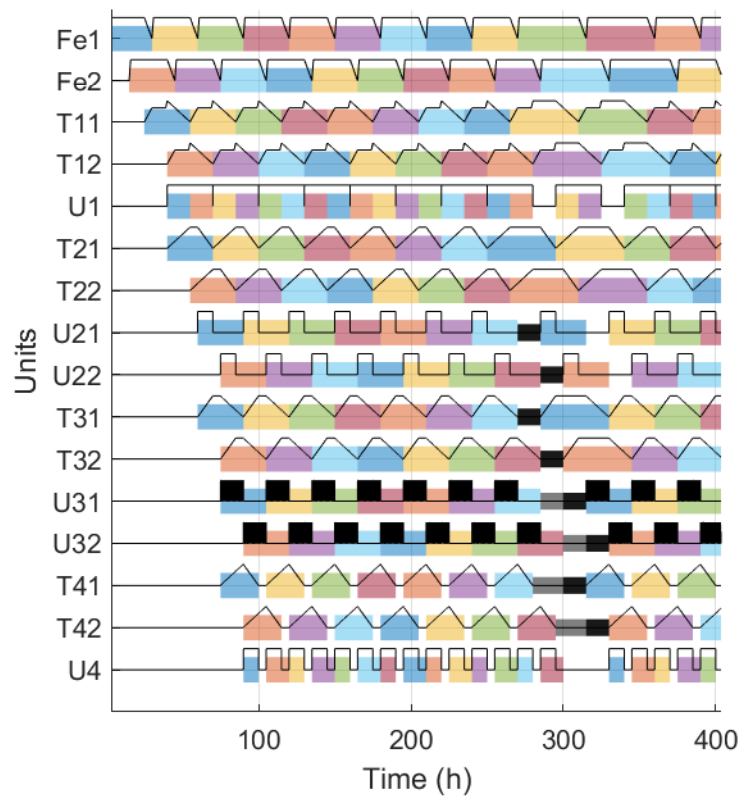
Revised Date: 27 March 2020

Accepted Date: 30 March 2020

Please cite this article as: Franz D. Böhner, Oscar A. Prado-Rubio, Jakob K. Huusom, Discrete-Continuous Dynamic Simulation of Plantwide Batch Process Systems in MATLAB, *Chemical Engineering Research and Design* (2020), doi: <https://doi.org/10.1016/j.cherd.2020.03.030>

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Highlights

- Modelling batch process systems w. shared resources in MATLAB/Simulink/StateFlow
- Visual cycle time analysis of different cleaning-in-place strategies
- Hybrid formalism which allows inclusion of continuous dynamics

Journal Pre-proof

1 Discrete-Continuous Dynamic Simulation of Plantwide 2 Batch Process Systems in MATLAB

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8 **Abstract**

Batch chemical and biochemical plants play an important role in the process industries. They are characterised by hybrid (continuous & discrete) dynamics as well as complex sequences and decision logic in the case of shared resources. This is challenging from a modelling and simulation perspective, both in terms of numerical algorithms as well as implementability and scalability/maintainability within software environments. In this work it is shown that it is possible to model complex plantwide batch processes at reasonably high performance, accuracy, and practicability in MATLAB/Simulink using the StateFlow toolbox. To this end, useful implementation guidelines are presented, and a complex example batch process is modelled. As focus lies on the implementability of complex batch control logic, the model is limited to mass balances. The simulation results are evaluated and carefully visualised, indicating the MATLAB's capabilities for analysis of such systems.

9 *Keywords:* Batch Process Systems, Process Modelling, Hybrid System,
10 MATLAB Simulink StateFlow, Cycle Time Analysis

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11 1. Introduction

12 Batch processes play an important role in the production of speciality
13 chemicals and pharmaceuticals (Croughan et al. (2015); Edgar (2004)).
14 Compared to continuous processes, they suffer from a number of
15 performance shortcomings (Teoh et al. (2016)). These include practical
16 limitations originating from operational (scheduling) complexity which can
17 lead to low equipment utilisation (Amaran et al. (2016)). Furthermore,
18 energy- as well as material integration of batch processes may require
19 advanced scheduling methods (Fernández et al. (2012)) or intermediate
20 storages (causing additional operational and/or capital expenditure).
21 This complexity may be seen as a good argument for the use of process
22 simulation in order to improve existing processes or design new ones (Foo &
23 Elyas (2017)). Simulation of batch process systems is challenging compared
24 to continuous processes as they ordinarily exhibit pronounced continuous
25 and discrete dynamics. Discrete events occur as elements of the sequential
26 batch control logic, whereas continuously integrated state space models are
27 usually needed to describe to underlying physiochemical processes.
28 Furthermore, Baldea & Harjunkoski (2014) point out the "inherently
29 non-stationary nature, the non-linearity of their models, and the dynamic
30 complexity that arises from the potential need to coordinate multiple units
31 and stages that operate in parallel".
32 In industry, continuous and discrete system domains of batch process
33 systems are ordinarily regarded in a separated approach: the scheduling
34 domain is optimised in discrete-event manner (or an alike abstraction which
35 focusses on material flows). Unit operation models may be based on
36 mechanistic or data-driven continuously integrated models; dynamic
37 plantwide phenomena are then omitted. Today, discrete-event simulation
38 may be regarded a standard method as a number of software vendors cater
39 to this market (for example ExtendSim[®], SchedulePro[®], INOSIM[®],
40 Simio[®], and AnyLogic[®]). By means of manual optimisation or

41 evolutionary algorithms they facilitate optimal equipment utilisation as well
42 as efficient conduction of engineering projects.

43 The concurrent simulation of continuous and discrete system-elements
44 (hybrid systems) on plantwide level is not yet a standard method. In
45 industry, in many cases, the merits of integrating these layers into one
46 simulation may not justify the added complexity. However, as Baldea &
47 Harjunkoski (2014) or Costandy et al. (2018) point out, the integration of
48 scheduling and control is an important area of research. This concerns
49 especially processes where time constants of the continuous subsystems are
50 so large that time-scale separation errors become relevant. Furthermore,
51 advances have been made in data-driven modelling. This may facilitate the
52 identification of complex reaction kinetics (Galvanauskas et al. (2018)),
53 thereby enabling the inclusion of such detailed effects also in plantwide
54 studies. In combination with abundant computation power - for instance
55 through cloud solutions - this renders the computation of rigorous hybrid
56 (here continuous-discrete) models feasible where it previously may not have
57 been. Finally, benchmark models are invaluable as a driver of research
58 progress and dissemination. To the best of the authors' knowledge, a
59 rigorous continuous-discrete model of a complex batch process system in an
60 academically accessible environment which enables the execution of
61 advanced methods is not documented in open literature. (A good example
62 for a continuous production plant is the re-implementation of the Tennessee
63 Eastman Benchmark Problem by Bathelt et al. (2015)).

64 Concluding, there are several reasons which render continuous-discrete
65 simulation of batch process systems an important topic of research. This
66 article specifically examines MATLAB Simulink's capabilities for simulating
67 complex plantwide batch process systems. It promotes using the
68 StateFlow[®] toolbox (MathWorks (2019b)) in order to graphically program
69 state charts (Harel (1987)). The use of StateFlow in modelling hybrid
70 systems is not new (Simeonova (2008); Sahbani & Pascal (2000)), but an

71 implementation which is comparable in terms of complexity with the
72 example plant presented in the article at hand is, to the best of the authors'
73 knowledge, not documented in open literature. Creating and maintaining
74 complex plantwide models is a challenging task and much focus in the
75 following is dedicated to handling this challenge in the MATLAB which,
76 unlike the discrete-event simulators listed in the above, has not been
77 specifically designed for this. On the other hand, it offers several
78 advantages such as flexibility as well as the inclusion of data analysis and
79 model building into one environment - which furthermore enables facile
80 implementation of advanced methods.

81 The article is structured as follows: section 2 gives an overview of
82 important classes of dynamic models as well as computational approaches
83 of simulating them. This is followed by section 3, which entails a step-wise
84 modelling framework. In the course of this, the article specifically discusses
85 how to address the structural challenges that arise when modelling complex
86 batch process systems in MATLAB/Simulink/StateFlow. Finally, the
87 simulation results of an exemplary process are presented in section 4. This
88 includes careful visualisation of the simulation results. The eligibility of the
89 software environment for these types of simulation studies, implementation
90 and computation challenges, and finally future work are discussed in section
91 5. Hereafter, the work is concluded.

92 2. Computation Approaches to Systems with Hybrid Dynamics

93 This section gives an overview of the discrete-event related part of the
94 dynamic systems classified in table 1. Here, GDEVS refers to the
95 (*generalized discrete event specification*) framework (Giambiasi & Carmona
96 (2006)). In hybrid systems, discrete events occur not only in the form of
97 (predictable) timers, but as a consequence of implicit, iterative algorithms
98 such as numerical solutions to systems of differential equations. This is
99 computationally challenging, and especially in optimisation studies with

100 large computational overhead it is thus essential that the modeller finds the
101 appropriate level of abstraction.

102 2.1. Discrete-Event Systems

103 In many instances it may be sufficient to neglect the continuous
104 elements which greatly reduces the computational burden. Examples of this
105 are discrete-event model based material flow analyses. They fit well the
106 ambition to *model with a purpose* (Daoutidis et al. (2018)): the evolution of
107 a tank level between *full* and *empty* is not usually essential for optimising a
108 plant schedule. As long as the flow itself is predictable enough to know the
109 points-in-time of full/empty, there is no information loss if a discrete-event
110 model is chosen.

111 These models have largely been developed within the operations research
112 community and are frequently used in discrete parts manufacturing, traffic
113 studies, or supply chain optimisation (Bangsow (2012)). There exist two
114 computation paradigms (Law & Kelton (2000)): the *next-event time*
115 *advance* and the *fixed-increment time advance* clock update. If occurrence
116 of the next event can be predicted in advance, next-event time advance is
117 favourable. A fixed-increment time advance algorithm not only takes
118 unnecessary steps, it also induces discretization error that scales with the
119 fundamental step size unless all events are strictly multiples of the chosen
120 sampling rate - this can be somewhat amended by using a variable-step
121 solver. Systems for which the clock value of the next event can be
122 computed in a straightforward manner are especially systems of timed
123 automata (Alur & Dill (1994)) as well as multi-rate timed automata (Alur
124 et al. (2000); Geist et al. (2008)). They allow discrete approximations of
125 systems comprising mass flows, storage tanks, and time-based operations in
126 batch production plants.

127 2.1.1. Discrete-Rate Simulation

128 Discrete-event models have been developed outside of the process
129 systems engineering domain, yet they are applied successfully for problems
130 arising within it (Petrides et al. (2014); Amaran et al. (2016)).
131 Discrete-rate simulation (DRS), available for instance within the
132 ExtendSim[®] environment, aims at amending some of the shortcomings of
133 classical discrete-event models such as the periods of stasis between events.
134 To this end, DRS allows states (inventories) to evolve on a continuous
135 linear envelope between events. Negligence to do this can, in some cases,
136 lead to accruing numerical error (Damiron & Krahl (2014)).

137 2.2. Hybrid Systems

138 Continuous behaviour is best described by systems of differential
139 equations. Analytical solutions to these systems are normally not
140 obtainable, thus a time-advance based on a list of next events is not
141 possible. The coexistence of continuous and discrete dynamics requires that
142 the solver is capable of handling both, and in the following two distinctions
143 are drawn which are likely to influence the choice of solver.

144 2.2.1. Hybrid Systems with Frequent Discrete Events

145 If system dynamics are largely dictated by discrete elements, one can
146 try to find local approximations of the continuous system trajectory in such
147 a way that they can be handled by a discrete-event solver. This is not
148 element of this article, and a large base of literature around the GDEVS
149 framework is available (Giambiasi & Carmona (2006); D'Abreu & Wainer
150 (2003); Giambiasi et al. (2001)). Furthermore, in the form of PowerDEVS
151 (Bergero & Kofman (2011)), a Simulink-like process simulator with
152 user-friendly graphic implementation features is available. These solvers are
153 necessary if discrete events occur at very high frequencies, for instance
154 during periods of *chattering*. Compared to the function evaluations needed
155 to describe continuous system behaviour, discrete events occur at low

156 frequencies in batch process systems. This frees the modeller from the need
157 to pursue such an approach.

158 *2.2.2. Hybrid Systems w. Infrequent Discrete Events*

159 Numerical integration is a standard method of chemical engineering to
160 solve systems of ordinary differential equations (ODE's). Also for stiff
161 systems, a variety of performant implicit solvers (in the case of MATLAB
162 for instance ODE15s, ODE23s) are available. Chemical engineers are
163 generally familiar with these methods and skilfully balance numerical error
164 with accuracy.

165 The question is then how to handle the discrete-event part of the system.
166 From a computational perspective, a split system approach is preferable
167 (Nutaro et al. (2012); Bouchhima et al. (2007); Clune et al. (2006)). In this
168 scheme, continuous and discrete system fractions are calculated
169 independently, and synchronisation only occurs when an event is triggered.
170 This is attractive both in terms of computational performance as well as
171 numerical error control during the numerical integration scheme.

172 However, especially in the case of complex systems with many elements,
173 links, and transitions, implementation of split models may be cumbersome.
174 Thus, one can choose to embed the discrete dynamics into the continuous
175 solver regime. A good balance between practicality and performance is
176 indicated by using variable step solvers with discrete-event detection. The
177 advantages in implementing and maintaining these models may outweigh
178 the disadvantages (accuracy, performance) as computation is fairly cheap.
179 Software environments capable of this are i.e. gPROMSs[®],
180 Modelica[®]/Dymola, or MATLAB/Simulink[®] (van Beek & Rooda (2000)).

181 *2.3. Hybrid Systems in MATLAB/Simulink*

182 The solver capabilities in place, several attributes render
183 MATLAB/Simulink attractive from a modellers point of view. Firstly, as
184 an environment apt both for data processing and modelling, it manages to

185 integrate two tasks which are ordinarily separated if dedicated process
186 simulators are chosen. Furthermore, through numerous libraries/toolboxes,
187 it allows facile implementation of advanced methods either within Simulink
188 flowsheets or in the embedding MATLAB environment.

189 However, neither the user interface nor currently available libraries are
190 designed for the implementation of complex batch process systems.
191 Notably, the SimEvents[®] (MathWorks (2019a); Gray (2007)) toolbox is
192 developed specifically for systems comprising discrete events and allows the
193 presence of continuous dynamics (Clune et al. (2006)). However, it is
194 optimised for systems consisting of queues and entities, which is not
195 practical for the implementation of sequential/parallel hybrid batch process
196 systems. This can be inferred from the predefined function blocks within
197 the toolbox (entity generators and sinks, queues, servers). While they are
198 useful elements of a high-level abstracted discrete-event study, they are not
199 convenient within the context of a hybrid simulation that includes
200 continuously evolving states. SimEvents expands Simulink by useful
201 elements connected to entity-management, which in a process systems
202 context is necessary for batch tracking. However, this can also be
203 implemented in the StateFlow framework with reasonable effort (shown in
204 section 3.2.7).

205 **3. A Framework for Modelling Batch Process Systems in** 206 **MATLAB/Simulink/StateFlow**

207 In the following, a stepwise procedure is introduced which separates the
208 modelling task into a series of sub-tasks. This is in general anticipated to be
209 of great help due to the complexity of the endeavour.

210 *3.1. Limitations of Continuously Solved Flow Charts*

211 In the chosen simulation approach, state charts have to be solved under a
212 continuous regime, which in MATLAB R2019b has the following implications:

- 213 • Library-links are disabled.
- 214 • No state transitions through event-broadcasting.
- 215 • Outputs cannot be written during state activity.

216 Furthermore, the absence of model libraries renders implementation tedious
217 and therefore error-prone. Event-broadcasting within a flowchart is
218 convenient in synchronising resources and callers (both of which there can
219 be multiple). However, dynamic updates of outputs (for instance set points
220 for manipulated variables) during state execution can usually be emulated
221 on root flowsheet level. Note also that, if the chart was a pure timed
222 automaton (all next future events are predictable at current event), it could
223 still be executed in event-based manner also within a continuously solved
224 flowsheet.

225 *3.2. Stepwise Model-Building Procedure*

226 In the following, the most important aspects of the model building
227 procedure in Simulink and StateFlow are elaborated. They are concisely
228 presented in figure 1, and each step is explicated in a dedicated sub-section.
229 It is not strictly necessary to follow this sequence, but there is a rather
230 natural order to it. In the proposed approach, a model has two layers: the
231 batch control system (a StateFlow state chart) and a physical process
232 counterpart. The latter is normally a system of differential-algebraic
233 equations, modelled using integrators or S-functions on root flowsheet level.
234 This is conceptually visualised in figure 2. The state chart layout is
235 representative of a StateFlow implementation.

236 *3.2.1. Step 1 - Model Configuration Parameters*

237 Aside from general solver properties (tolerances, algorithm, etc.), the
238 number of batches to be processed during a campaign is best specified in
239 advance. A sufficiently long simulation horizon to process all batches should

240 be chosen; the simulation can be terminated prematurely when the last
241 batch has been processed on the most-downstream units (and all machines
242 have returned to idle state after finishing the last re-initialisation).
243 Simulating over such large time horizons may require controlling major
244 integration step size as the default step size is large if left on automatic
245 selection. In general, despite of the use of StateFlow under a continuous
246 solver regime, cases are experienced where events are not properly detected
247 if input values change rapidly compared to step size. This will usually be
248 identifiable by implementing a series of simulation integrity checks (section
249 [3.4](#)). Therefore, adjusting model configuration parameters (step 1) is
250 iterative by nature. This is especially so if no units with inherent step size
251 requirements are installed (for instance pulse- or sine sources). Choosing
252 the second-largest step size (by order of decimal place) which leads to exact
253 solutions has shown to be a robust approach with good computational
254 performance. The fact that numerical error in the discrete-event system
255 part may occur is undesirable and the modeller needs to be alert.

256 *3.2.2. Step 2 - Define Structure of Process System*

257 Liquid or gaseous material must at all times be contained in a tank or
258 an equivalent storage unit. Hold-ups of flow processing units (centrifuges,
259 filters, etc.) are likely negligible. Systems can however be modelled to such
260 a high degree of fidelity if that is required. In the case of solids that can be
261 stored more flexibly, the requirement for a containing unit is relaxed. It is no
262 problem to extend the model by storages with room for more than one (solid)
263 batch, but this has not been implemented in this example. In the same way,
264 it is not a problem to combine batches in one tank or split a batch in two. It
265 is intuitive to choose a distributed modelling approach for the continuous part
266 of the system. Each physical entity that stores material (buffer tanks, unit
267 operations, etc.) is represented by an integrator or S-function. A second class
268 of physical units are those that predominantly process material downstream
269 (or recycle it) - continuously operating units. These can often be understood

270 as resources needed by the tank units, which also require a free recipient
271 tank before material can be sent downstream. Flow processing units can be
272 modelled using arbitrary (for instance algebraic) function blocks. Naturally,
273 also a complex dynamic model - embedded within for example an S-function
274 - can be implemented.

275 3.2.3. Step 3 - Batch Control System

276 Deciding whether or not to decompose the batch control system into
277 separate StateFlow state charts is less straightforward. Separate state
278 charts are generally more intuitive to understand as they comply well with
279 the concept of recipe-driven unit procedures. As the continuous solver
280 regime forbids the use of library functions within the state chart
281 environment, this furthermore allows mimicking object-oriented
282 programming: local variable names can be re-used within separate state
283 charts, which can therefore be copied easily. On the other hand, a
284 separation leads to more complex signal routing on root flowsheet level:
285 each time a variable is passed between charts, this requires that a link
286 (graphic or virtual) is drawn. Finally, the decomposition into multiple
287 charts led to stability issues in MATLAB R2019b and previous versions.
288 These issues can be circumvented by choosing a fixed-step solver, which is
289 however undesirable due to performance and accuracy set-backs (section 2).
290 A centralised implementation (one superordinate flow chart) requires that
291 the embedding chart is solved under *parallel* (AND) decomposition. It
292 contains an embedded sub-chart for each unit and storage tank present in
293 the system. These are all executed at the same time and initialised as *Idle*.
294 They represent the actual machines which can only ever be in one state and
295 must themselves be solved under *exclusive* (OR) decomposition.
296 *Virtual units*: it is possible to model resources within the control system
297 that have no representation in the actual flowsheet. This might be
298 convenient if they have no inherent dynamics and there is no interrelation
299 with the process other than an effect on the schedule. Examples of this are

300 cleaning-in-place stations or operators.

301 3.2.4. Step 4 - Interface: Batch Control System - Process

302 As indicated in figure [1](#), the layout of this interface is closely
303 interrelated with the two consecutive steps (resource handling & material
304 routing). Firstly, the interface needs to contain ports which allow passing
305 measurements from the process to the control system. In a scenario limited
306 to mass flows, this specifically concerns tank volumes/levels, but it is easy
307 to extrapolate to temperatures, pH, or any variable obtained from a direct
308 or inferential measurement. Secondly, the bi-directional interface needs to
309 be able to pass command signals from the batch control system to the
310 physical process. From a simulation perspective, these signals can be
311 directly passed to the units. Yet, if the modeller chooses to do so, it is easy
312 to implement a regulatory control layer in-between. This also enables the
313 introduction of *implementation errors* in the control loop (useful for
314 diagnostic studies), and extends the model by dynamics from feedback
315 components in the lower layer. In practical terms, the interface depends on
316 the structure of inputs and outputs of state chart and (unit operation)
317 sub-systems. As a centralised state chart implies that numerous variables
318 are passed on, these should conform to an intuitive and consistent
319 nomenclature and array sequence (channel 1 - inflow, channel 2 - outflow,
320 ...).

321 Beyond the signals on unit operation level, the control system needs to be
322 able to implement resource handling / material routing on flowsheet level.
323 Therefore, a further variable which essentially emulates valve positions in
324 the piping system is necessary.

325 3.2.5. Step 5 - Resource Handling in Control System

326 If a plant layout is fixed (each upstream tank feeds via a standard unit
327 to a fixed downstream tank), an implementations in StateFlow is trivial
328 and this step is reduced to checking whether the statically assigned units

329 are *Idle*. However, in modelling flexible batch plants, resource handling is
330 one of the key challenges. It is exacerbated by the limitation that state
331 charts cannot broadcast events under a continuous solver regime (section
332 3.1). This precludes the immediate transition of a resource unit from *Idle*
333 to *Busy* upon being claimed by an caller.

334 A workaround is possible as all sub-charts within a superordinate chart can
335 write to- and read from the local state chart workspace. Note that, also as
336 a consequence of continuous solving, library functions are disabled, and the
337 list of variables can quickly grow very long. Therefore, thoughtful and
338 consistent naming of variables is essential.

339 Resource handling firstly entails the check for a free process path (often
340 both, a processing unit and a recipient tank). The check has to cover all
341 related downstream units; introducing a dedicated state for each path
342 (state *toTank1viaUnit2*, ...) allows choosing preferred recipient units by
343 assigning an order to the state transitions. The moment an upstream unit
344 starts sending material downstream, *processing resource* and *recipient tank*
345 must no longer be called from a further upstream unit. This is best
346 controlled by the origin tank, and here the state chart workspace comes
347 into play. As events cannot be broadcast to the resources (causing them to
348 leave *Idle*), a *setter variable* in the shared workspace accounts for the
349 utilisation state of the unit. The upstream caller immediately assigns this
350 variable a new value when a resource is called. It should now be evident
351 that upstream units check the state of this variable to find a feasible path.
352 As resources might be busy in a procedure which is not linked to an
353 upstream caller (re-initialisation, cleaning, maintenance), a check for *Idle* is
354 still necessary. Each resource can reset the setter variable once the
355 procedure is complete and the link to the caller broken.

356 3.2.6. Step 6 - Material Routing on Flowsheet level

357 It is crucial that the flow of each unit is specified only by one control
358 system caller. Furthermore, every outflow must eventually turn into the

359 inflow to another storage unit. In the case of systems of multiple sources
360 and destinations, coordination of these mass flows is necessary. The most
361 intuitive way to solve this is by using Simulink selector blocks which contain
362 the necessary functionality: only one input can be passed through at a time
363 (indicated by bold lines in figure. Beyond this, only inflows from external
364 sources are allowed in downstream tanks. (Unlike in Modelica, in Simulink
365 the modeller is not forced to model connectors in a way that promotes mass
366 balance consistency.) The flow resource is the counterpart to a connector,
367 and in this way, the material balance throughout the plant is closed.
368 The setter variable connected to a resource (introduced in the previous
369 section) is helpful in the material routing problem: it is efficient to use it
370 not only to hold (*Idle/Busy*) information, but to contain the ID of the
371 upstream caller. If callers are enumerated regularly, these ordinals can be
372 used to control the path through the selector blocks. If no caller is
373 specified, the variable is to hold zero. This index is chosen as the default
374 feed-through of the selector blocks - if this signal is also zero, it does not
375 affect the mass balance. (Disabling sub-systems of tanks and other
376 resources during inactivity creates redundancy.)

377 3.2.7. Step 7 - Implement Batch Tracking System

378 The functionality to track batches through the system is not required
379 in order to be able to execute the simulation. However, it is important in
380 posterior validation as well as evaluation of a simulation study. A batch
381 ID can be created either in the queue or in the most-upstream unit, in the
382 simplest case it counts up incrementally, which is easy in StateFlow and
383 only requires a further local variable. As a receiver knows by which unit
384 it was called, it can take over the batch ID from upstream units. If there
385 is a dedicated state in each unit for each caller-resource / sender-resource
386 combination (section [3.2.5](#)), this is implementable with ease.
387 Not only the batch number is required to keep track of all statistics, the
388 machine states (*Filling, Waiting for ..., ...*) need to be logged as well. Also

389 here, consistent naming is crucial to render the system as understandable as
390 possible. In this work, machine steps have been classified according to the
391 following keys:

- 392 • 0 idle
- 393 • 1,2,3,... standard operations (nominal processing)
- 394 • 100,200,300,... waiting for resources / recipients during processing
- 395 • -1 re-initialisation
- 396 • -10 CIP called
- 397 • -11 CIP in progress

398 During nominal operation, the first cypher counts up continuously: an
399 exemplary sequence reads 1, 2, 300, 4, 500, 6, -1 for a unit with four
400 operations (1, 2, 4, 6), two waiting steps (300, 500), and a re-initialisation
401 step (-1). There is no standard number associated with a certain type of
402 step (filling/emptying), and it is entirely up to the modeller to find an
403 appropriate enumeration. Similarly, the assignment of negative values for
404 re-initialisations and CIPs is arbitrary. Here it is chosen such that it
405 facilitates selective plotting/colouring schemes.

406 3.3. *Re-initialisation vs. Cleaning-In-Place (CIP)*

407 With the above, the basics for putting together a functional batch
408 process system in Simulink using StateFlow state charts are in place.
409 However, several functionalities necessary for realistic modelling have not
410 yet been introduced. These are i.e. equipment re-initialisations and CIPs.
411 The distinction is drawn as CIPs or sterilisation-in-place (SIPs) are
412 understood as plantwide issues which require coordination with other units
413 and a CIP system (resource), whereas unit operation re-initialisations are
414 local. Inclusion of a virtual CIP station (free/busy) in the batch control

415 system is sufficient, naturally a *physical entity* can be implemented on root
416 flowsheet level if this is desirable. Either process may require operator
417 attendance (resource); the operator model (busy/free/activity) can be
418 implemented in the same way as a virtual CIP station. A re-initialisation or
419 CIP may be called after each batch, after a certain time, after a certain
420 number of batches, after a certain event-occurrence on a batch, or after a
421 transient variable for some unit operation (for instance fouling) crosses a
422 threshold.

423 3.4. Validation

424 As indicated by [Tiwari \(2002\)](#), especially for large systems with complex
425 input patterns it is a challenging task to verify whether a state machine
426 reacts to arbitrary input patterns in the desired way. To this end, formal
427 verification methods exist that are not element of this work. However,
428 validation of complex batch campaigns is an issue that needs to be
429 addressed. The identification of faulty sequences on a unit operation is not
430 a problem, as visual verification for several scenarios can, with relatively
431 high certainty, confirm that the sequence is implemented properly.
432 (Furthermore, implementing sequences on unit operation level is relatively
433 straightforward.)

434 Resource handling and material routing are substantially more error-prone.
435 It would greatly upset the fidelity of a simulation if a batch could be lost or
436 created 'out of nowhere' in the middle of the downstream line. (Or, in the
437 worst case, both - which renders detection difficult.) Keeping track of the
438 total number of processed batches allows to deduce whether simulation
439 integrity is maintained or not. Visual or automatic checking of cycle times
440 and volume profiles under different solver options gives the modeller a quick
441 feeling for both implementation and numerical issues.

442 4. Implementation of an Example Plant

443 In the following, the implementation of a reference batch process system
444 is presented. It is inspired by the industrial case study documented in
445 [Bähler & Huusom \(2019\)](#). Beyond modelling, focus is put on posterior
446 graphic evaluation to give the reader an intuitive understanding of the type
447 of process which has been simulated. In terms of operational complexity
448 (i.e. dynamic selection of units and cleaning operations) the model does not
449 stand back from batch or hybrid processes documented in open literature
450 ([Montes et al. \(2018\)](#); [Alskehli et al. \(2010\)](#); [Monroy & Vallejo \(2013\)](#);
451 [Sharda & Bury \(2010\)](#); [Toumi et al. \(2010\)](#); [Noguera & Watson \(2004\)](#)).
452 The generic example line is fed by two fermenters; before each set of unit
453 operations, two holding tank are installed. Some process steps contain
454 parallel machines, others only one. A schematic overview is given in figure
455 [3](#).

456 4.1. Model Specifications

457 A description of the operational procedures in terms of constraints and
458 rates is given in table [2](#). A simple campaign is visualised in Gantt chart
459 notation in figure [4](#). Here, a campaign of four batches (colour-code) is
460 shown. Scaled volumes and flow profiles are plotted for the sake of
461 understandability. A re-initialisation occurs when a unit is still *coloured*
462 due to attribution to a batch while no flows are processed (units U21/22,
463 U31/32, and U4). Flows on units U31/32 are chosen to alternate frequently
464 to show the possibility. Rapid changes in flow rates may require step size
465 control (section [3.2.1](#)) due to numerical error.

466 Cycle times on the machines are designed such that, normalised for the
467 number of available units, all process steps take 15 hours per batch. That is
468 with one exception: due to the irregular re-initialisation schedule on unit 4
469 (after every second batch), there is minor theoretical overcapacity. This
470 stems from the instances in which the machine is idle while no

471 re-initialisation occurs, as the timeslot needs to be reserved if a fixed
472 schedule is to be implemented.

473 4.1.1. Example CIP Procedure

474 To study complex schedules arising from cleaning-in-place events, the
475 model is extended by CIP routines. This firstly requires the introduction of
476 a CIP station in the batch control system (here reduced to a virtual
477 resource, section 3.2.3). CIP stations are often shared between production
478 lines and may block more than one machine at a time, for instance when a
479 tank is needed in order to CIP a unit. Therefore, all machines that are
480 subject to a CIP need to be extended by the related states, these are i.e.
481 *CIP called* and *CIP in progress*.

482 The scenario is designed such that a CIP covers units 21/22 up to tanks
483 41/42, as indicated in figure 3. It is called every 8.75 days, thus fitting
484 exactly into the schedule. It follows a rigid procedure which can be seen in
485 figure 5. In the standard CIP sequence, firstly unit 21 and tank 31 are
486 cleaned congruently. Upon completion, unit 31 and tank 41 are blocked to
487 enable a consecutive CIP (grey bar). Once unit 22 and tank 32 have
488 processed the last batch, they proceed to active CIP. Blocking the
489 downstream units from further processing guarantees a coherent CIP
490 barrier between the pre- and post-CIP batches. When the CIP on unit 22
491 and tank 32 is completed, unit 31 and tank 41 can proceed to active CIP,
492 and as in the above, unit 32 and tank 42 are blocked from processing a
493 batch to prevent cross-contamination until they are cleaned in the final CIP
494 routine. Each cleaning of a unit/tank group lasts exactly 15 hours - the
495 constraining cycle time in the system.

496 A second procedure follows a different pattern: after cleaning U21/T31,
497 U31/T41 are subjected to a CIP. Consecutively U22/T32, and finally
498 U32/T42 are cleaned. This is shown in figure 6, and visual assessment
499 reveals that less waiting is experienced in this scenario.

500 *4.1.2. Effect of CIPs on Equipment Efficiencies*

501 It can be seen that the first introduced CIP procedure in figure 5 leads
 502 to a significant amount of blockage due to units being taken out of
 503 operation in anticipation of a CIP. As the CIP duration is designed such
 504 that it actually fits into the schedule, this is suboptimal and leads to long
 505 cycle times induced by the step *waiting for CIP* (step -10), visualised in
 506 figure 7-a. Cycle times of the second (improved) procedure are presented in
 507 figure 7-b. They lie notably below those in the previous CIP design and
 508 lead to a 5.4% capacity increase.

509 The schedule indicated in figure 6 (dashed rectangle) exhibits an interesting
 510 property, namely the coincidental starting and finishing of the filling
 511 procedure in tanks 41/42. In the designed case it does not matter in which
 512 order the tanks are processed on unit 4, as it leads to equal waiting periods.
 513 It is however a good example of complex scheduling decisions which are not
 514 trivial to make without support through technological tools, as it would
 515 require the proper course of action if capacities were leveraged.

516 *4.2. Behaviour Under Stochastic Uncertainty*

Randomised waiting can easily be added in within the control system by
 introducing a random timer. This is representative of manual control, as
 operators may react delayed. Beyond that, the physical system on Simulink
 flowsheet level can easily be extended by random effects. As a simple
 example, in the following the flow rates on U21/U22 unit are subjected to
 Gaussian noise (flow rate values are kept constant during a cycle). To this
 end, the inlet flow rate of $1.5m^3/h$ is superposed a random term $\Delta_{F,i}$ with

$$\text{var}(\Delta_{F,U21}) = 0.05 \quad (1)$$

$$\text{var}(\Delta_{F,U21}) = 0.075 \quad (2)$$

517 As equipment capacities in the process are relatively even, system
 518 performance does not benefit from short cycles, but prolongations

519 propagate up- and downstream. An overview of the affected step durations
520 on tanks 21/22 is shown in figure 8, and it becomes evident that variability
521 is not only experienced in the affected material transfer step, but also
522 periods of idle time are experienced due to short cycles and upstream
523 delays. The according cycle times are shown in figure 9.

524 4.3. Computational Performance

525 Due to the great number of additional function evaluations under a
526 continuous solver regime - compared to a pure discrete-event system -
527 performance differences are present. The overall computation workload is
528 strongly linked to the differential equation solver, the nature of the
529 continuous system, and the allowed for error tolerances. A rigorous hybrid
530 simulation of a campaign of several batches may thus result in substantial
531 execution times and should be kept in mind.

532 In the example simulation, the system is reduced to piecewise linear mass
533 balances between time- and state discrete events. The system is solved with
534 ODE15s on an Intel® Core™ i5-5300U CPU which is rated at 2.3 GHz.
535 Execution time scales linearly with the number of batches in a campaign
536 and the Simulink flowsheet for a duration of 500 batches executes in less
537 than 25 seconds.

538 In this calculation, a maximum step size of 0.1 - which in the given system
539 corresponds to hours - and the standard absolute and relative ODE15 error
540 tolerances are chosen. The maximum step size does not equate incurred
541 error on event detection; the identification of these events is accurate and in
542 the case of intrinsic timers exact. However, if this tolerance is left
543 unchecked, the solver tends to miss chains of events entirely as the linearity
544 of the continuous system may trigger excessively large integration step
545 sizes. These can result in missed zero-crossings of the event detection
546 system. Increasing maximum step size to 1 (hour) reduces the execution
547 time to under 20 seconds; all events are still identified properly.

548 Performance after the inclusion of complex continuous dynamics remains to

549 be investigated, but at least it is indicated that the execution of the
550 discrete part of the system can be included in the holistic modelling
551 approach without inhibiting performance drastically.

552 5. Discussion

553 It can be concluded that the MATLAB/Simulink/StateFlow
554 environment is apt for modelling and simulating batch process systems.
555 This is expected, not least due to the fact that it has been used for
556 applications of this kind before. However, the implementation which has
557 been presented in the article at hand surpasses them in complexity, and
558 guidelines have been introduced which aid in structuring a complex model
559 building process. Overall, there are only few environments which can
560 handle true continuous-discrete models, especially when it comes to systems
561 with numerous elements and complex sequential procedures (such as batch
562 process plants).

563 The Simulink/StateFlow environment allows facile study of the interplay
564 between continuous and discrete systems. An exemplary phenomenon of
565 interest would be the effect of proportional-integral controller tuning on
566 time-scale separation error. Another example would be the quantification of
567 the gains from being able to terminate a fermentation subject to biological
568 variability based on process analytical technology rather than a fixed
569 schedule. In general, if dynamic phenomena connected to product quality
570 or yield are in need of quantification, this calls for an environment which
571 can handle continuous system elements (Costandy et al. (2018)). While the
572 models in Simulink/StateFlow are not accessible to mixed-integer solvers,
573 black-box optimisation algorithms can be tried. Furthermore, the models
574 can be used for the sake of validating an abstracted optimisation model
575 (Vieira et al. (2019)).

576 The proposed framework allows the generation of hybrid data sets based on
577 mechanistic models which resemble those of real production sites to a very

578 high degree. Here, the Simulink flowsheet environment gives the modeller
579 intuitive control of the inputs and thus the occurring effects. The modelled
580 behaviour can exceed mere random uncertainties, which is expectedly an
581 Achilles' heel of many machine learning algorithms. Therefore, this
582 framework might be seen as a first steps toward creating a sandbox
583 environment for facile testing and validation of data-driven algorithms
584 before they are tried in real production environments. (Here, typically a
585 large number of unknowns and uncertainties are beyond the analyst's
586 control). The value of accepted benchmark models such as the Tennessee
587 Eastman Proces (Downs & Vogel (1993)), or the Benchmark simulation
588 model no 2 (Jeppsson et al. (2007)) in disseminating maturity and aptitude
589 of technologies between academic but also industrial researchers has been
590 pointed out many times (see for instance Huusom (2015); Downs (2012)).
591 Unfortunately, numerical accuracy of the simulations needs to be asserted
592 through integrity checks which are likely to require some manual
593 evaluation. On the other hand, despite of Simulink's well understood
594 capabilities for solving continuous systems (Klee & Allen (2016)), it is not
595 unusual that simulation accuracy and performance are balanced by means
596 of iterative tuning. Therefore, this should not be considered a disadvantage
597 compared to other software environments. Simulation studies of campaigns
598 consisting of several batches are likely to require lengthy computation,
599 therefore this approach is not apt for real-time hard tasks.
600 While the chosen environment allows modelling batch process systems of
601 some complexity, limitations arise which one needs to be aware of.
602 Generally, it is not advisable to model complex

- 603 • Multi-purpose plants (plants without a fixed topology)
- 604 • Multi-product plants with severely different recipes (differing not only
605 in parameter values, but recipe sequences)

606 To some extent, this is a consequence of the restriction which arise in

607 continuously solved state charts (such as limitations in object-oriented
608 modelling practice and event broadcasting). Therefore, the model-building
609 process is likely too tedious and a recipe-based discrete-event simulator is a
610 much more appropriate environment. Not least, it is unlikely that detailed
611 reaction kinetics and unit operation models for a wide variety of products
612 are available.

613 On the other hand, for plants with fixed layouts and moderate product
614 diversity (or dedicated to one product) it is possible to build models in a
615 straightforward, graphically supported way. Here, the benefits of the
616 integrated MATLAB environment (data pre-analysis and parameter
617 identification, modelling & simulation, posterior data analysis and
618 optimisation) can be exploited - while offering facile inclusion and study of
619 continuous effects. Furthermore, simple manual scheduling studies
620 (Georgiadis et al. (2019)) which are ordinarily conducted in discrete-event
621 simulators can be executed effectively. MATLAB could be assessed with
622 respect to its abilities for pure discrete-event system studies of complex
623 batch process systems. While the modelling effort is still going to exceed
624 that in simulators dedicated to the cause, recipes can be implemented in a
625 straightforward way using StateFlow, especially as the limitations from
626 section 3.1 are mediated if a pure discrete-event solver is chosen.
627 Furthermore, the SimEvents[®] toolbox offers (strongly abstracted) standard
628 blocks which might be useful for such models. MATLAB's integrated
629 functionalities and widespread availability (i.e. due to its academic
630 licensing scheme) would enable effective method development, for instance
631 related to automatic derivation or validation of discrete-event models based
632 on batch process data.

633 6. Conclusion

634 In this work, applied guidelines have been presented that support
635 constructing sequential/parallel hybrid batch process system models in

636 MATLAB. An example plant has been simulated, and the capabilities for
637 posterior data visualisation and analysis have been shown. Model-building
638 in MATLAB entails some challenges which arise on the one hand from the
639 lack of standardised functionalities, and secondly from several limitations in
640 StateFlow as a consequence of a continuous solver regime. These difficulties
641 render the environment unattractive for industrial applicants who need
642 quickly-implementable solutions. Furthermore, it is inapt for systems with
643 high combinatorial complexity; still, it is shown that the simulation
644 environment allows the creation of holistic, non-linear, continuous-discrete
645 plantwide models of reasonably complex systems. Data sets can be
646 generated which closely resemble those of real batch process systems - with
647 full and intuitive control of the modelled phenomena and especially
648 disturbances. In the future, an implementation of a batch process system
649 benchmark model in MATLAB would enable easy access throughout the
650 academic community as well as facile testing and development of new
651 methods.

652 Acknowledgements

653 The authors would like to acknowledge the Technical University of
654 Denmark (DTU) and BIOPRO2. The project received financial support
655 from Innovation Fund Denmark through the BIOPRO2 strategic research
656 center (Grant number 4105-00020B).

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Type	Linear	Non-linear
Continuous Systems	<i>Continuous LTI models (Discrete LTI approximations*)</i>	<i>Differential (P)D(A)E Systems</i>
Discrete-Event Systems	<i>Multi-Rate Timed Automata</i>	<i>Automata, Petri Nets</i>
	<i>Frequent Events: GDEVS Formalism</i>	
Hybrid Systems	<i>Scarce Events: Discrete-Rate Simulation**</i>	<i>Scarce Events: (Complete Batch Process Systems)</i>

Table (1) Dynamic system types and common mathematical model expressions.

*Discrete computation of linear time-invariant (LTI) systems is trivial as the sampling rate is constant.

**Continuous evolution of volumes between events is considered in this otherwise discrete modelling framework (section [2.1.1](#)).

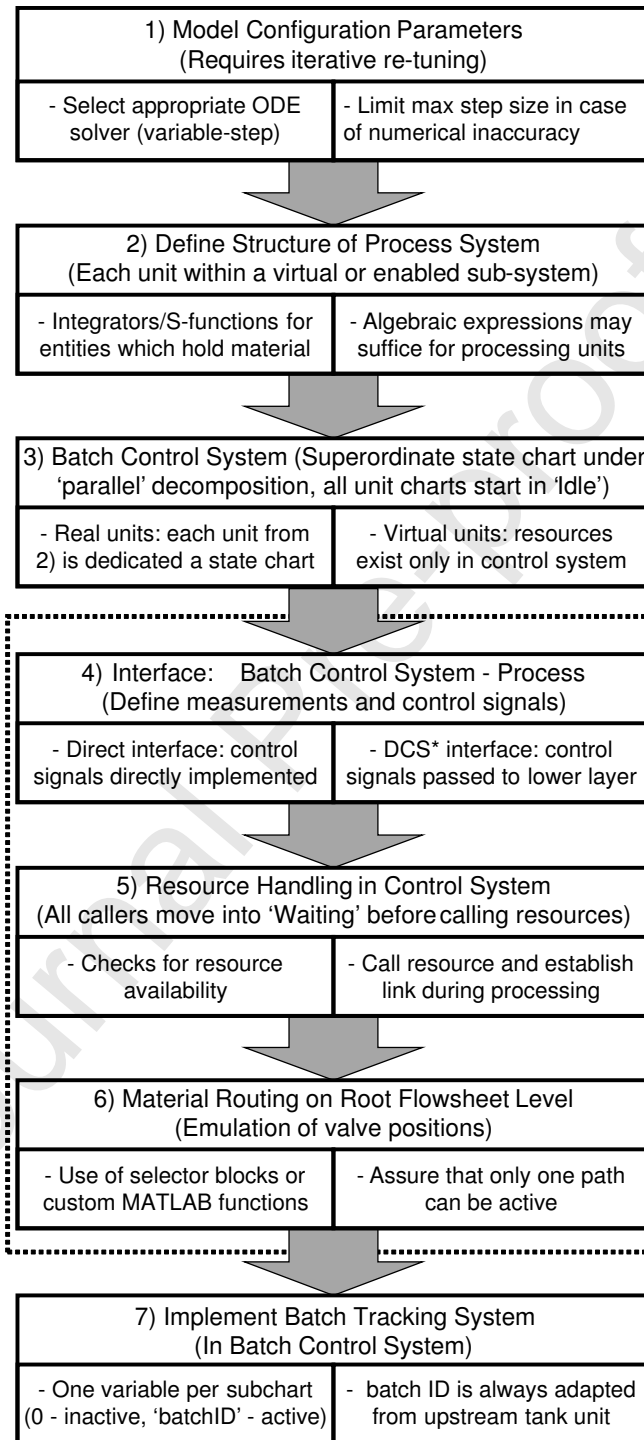


Figure (1) List of the most important steps concerned with building a batch process system model in MATLAB/Simulink/Stateflow. The steps are delineated in detail in the subsections within section 3.2

(* DCS: Distributed Control System)

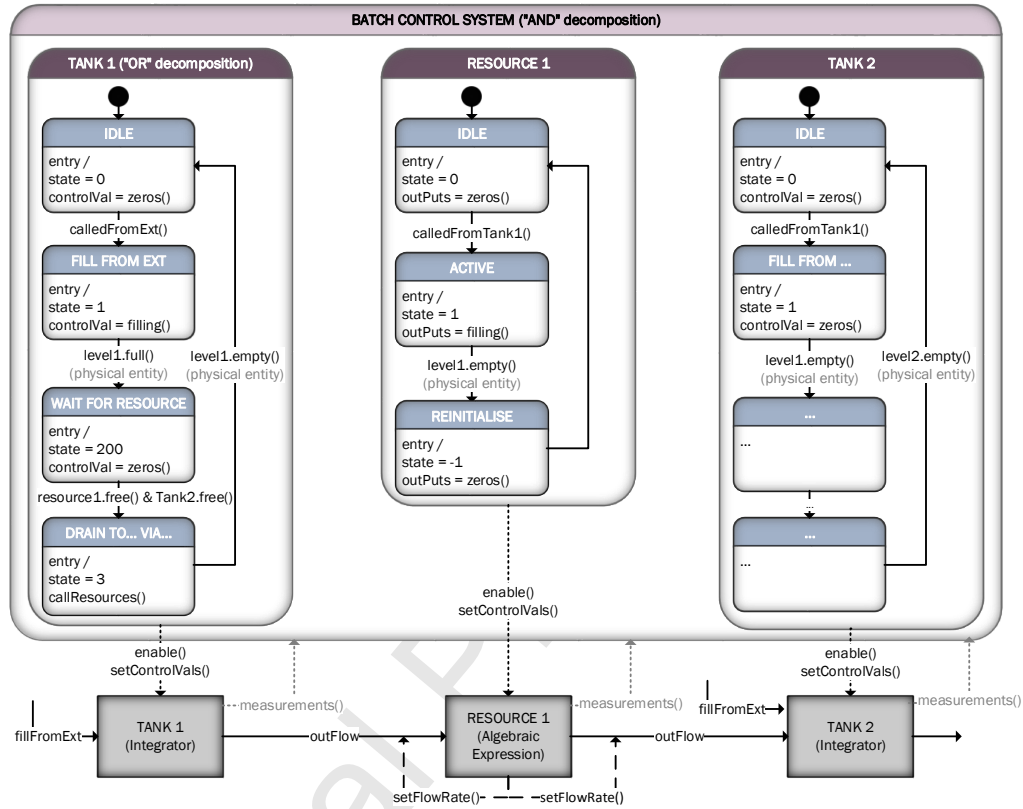


Figure (2) Exemplary architecture if one superstate with parallel decomposition of sub-states (machine states) and continuous elements on root flowsheet level is chosen.

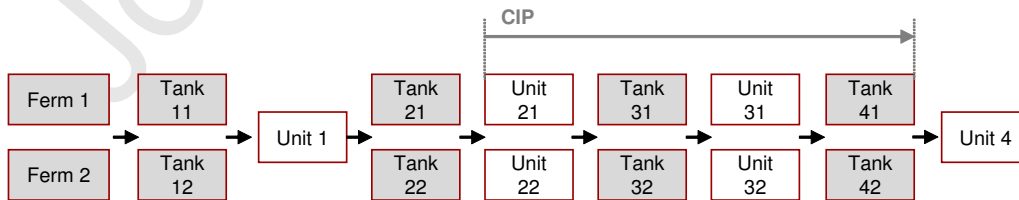


Figure (3) Overview of units in production line to be modelled.

Unit	Constraint			Dur. (h)
	Value	Rate	ID	
Steps (excl. idle)				
Ferm 1/2				
Fill from ext.	10 m ³	10 m ³ /h	1	1
Ferment	24 h		2	24
Wait for R.			300	
Drain	0 m ³	2 m ³ /h	4	5
Tank 11/12				
Fill	10 m ³	2 m ³ /h	1	5
Hold	9 h		2	9
Fill from ext.	+5 m ^{3*}	5 m ³ /h	3	1
Wait for R.			400	
Drain (via Unit 1)	0 m ³	1 m ³ /h	5	15
Unit 1				
Processing	15 m ³	1 m ³ /h	1	15
Tank 21/22				
Fill	15 m ³	1 m ³ /h	1	15
Hold	5 h		2	5
Wait for R.			300	
Drain (via U21/22)	0 m ³	1.5 m ³ /h	4	10
Unit 21/22				
Processing	15 m ³	1.5 m ³ /h	1	10
Reinitialise	15 h		-1	20
Tank 31/32				
Fill	15 m ³	1.5 m ³ /h	1	10
Hold	5 h		2	5
Wait for R.			300	
Drain (via U31/32)	0 m ³ /h	1 m ³ /h	4	15
Unit 31/32				
Processing	15 m ³	1 m ³ /h	1	15
Reinitialise	15 h			15
Tank 41/42				
Fill	15 m ³	1 m ³ /h	1	15
Wait for R.			200	
Drain (via Unit 4)	0 m ³	1.5 m ³ /h	3	10
Unit 4				
Processing	15 m ³	1.5 m ³ /h	1	10
Reinitialise	5 h**		-1	5

*Amount of material added relative to current fill level.

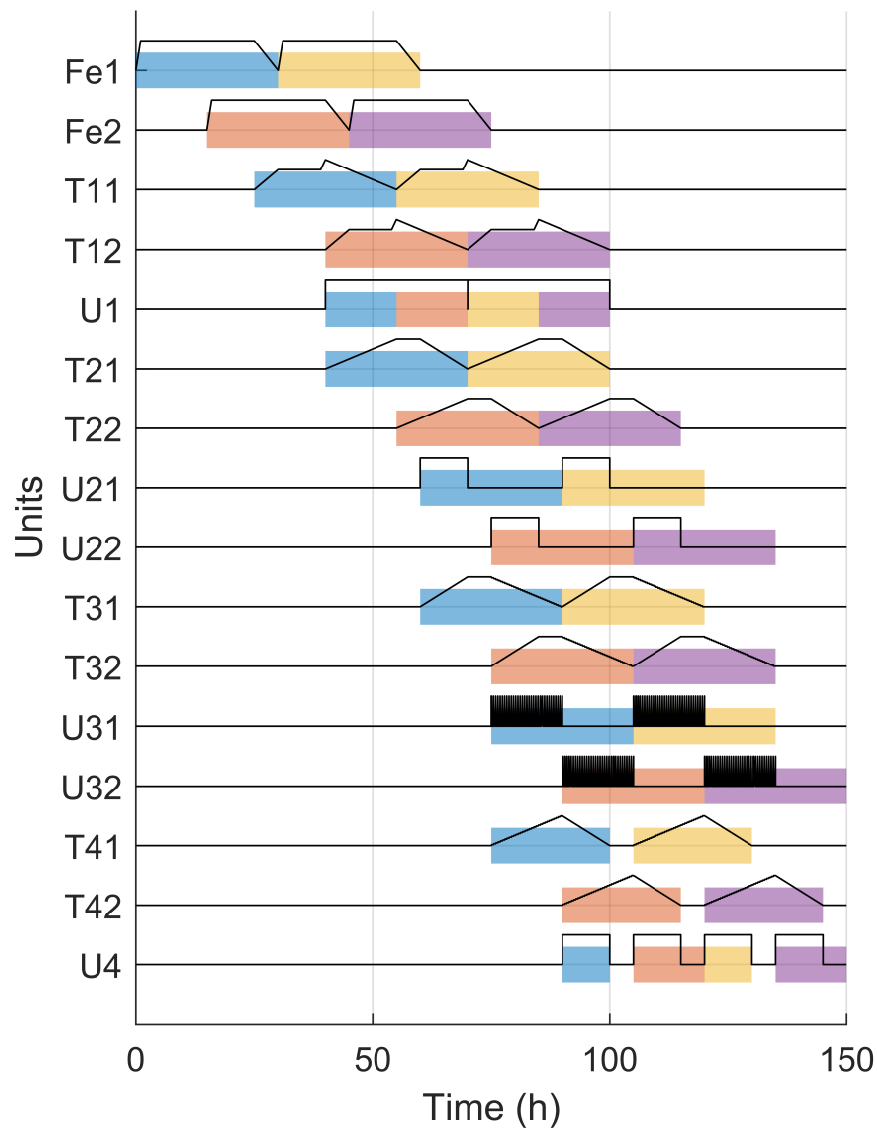


Figure (4) Exemplary campaign of four batches (colour-code).

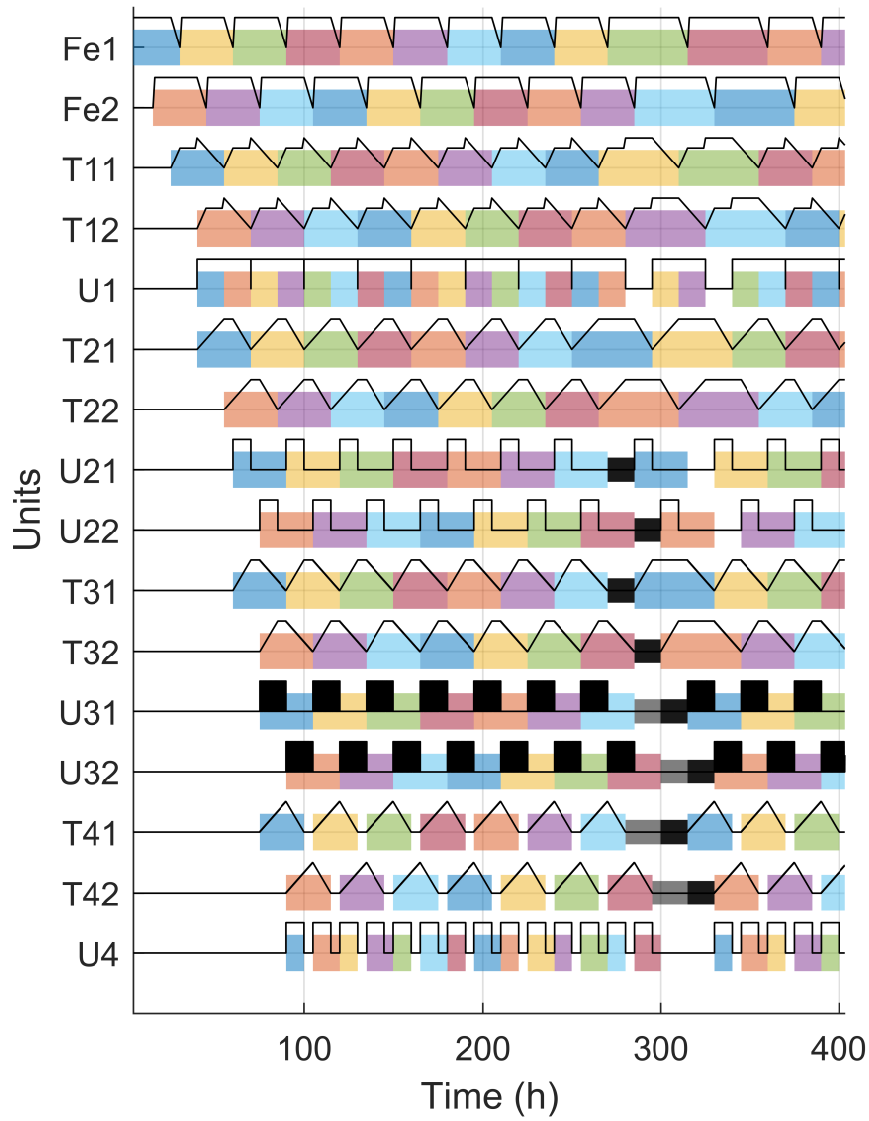


Figure (5) CIPs on units 21/22 & 31/32, tanks 31/32 & 41/42. Gray bar: unit blocked, CIP system busy. Black centred bar: CIP.

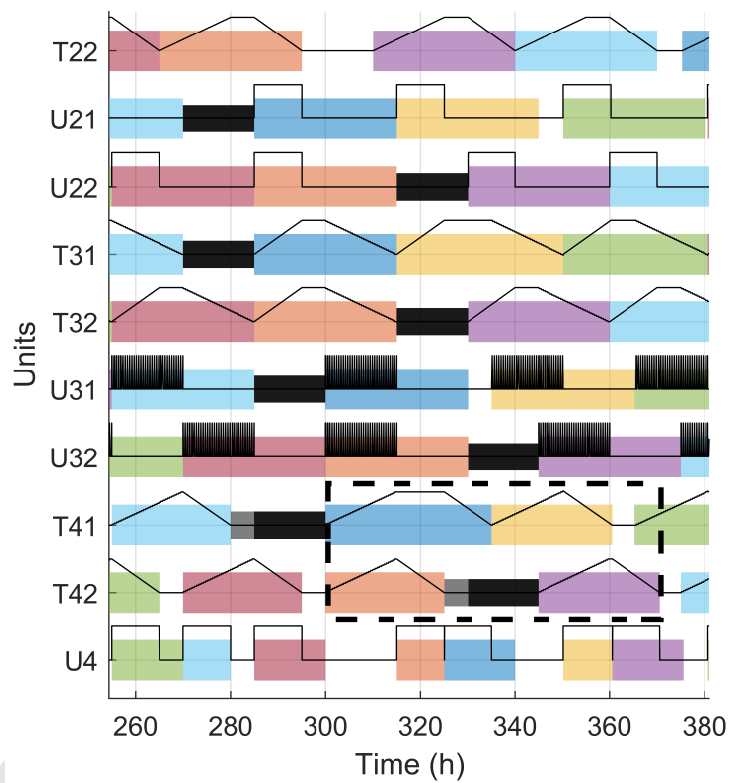
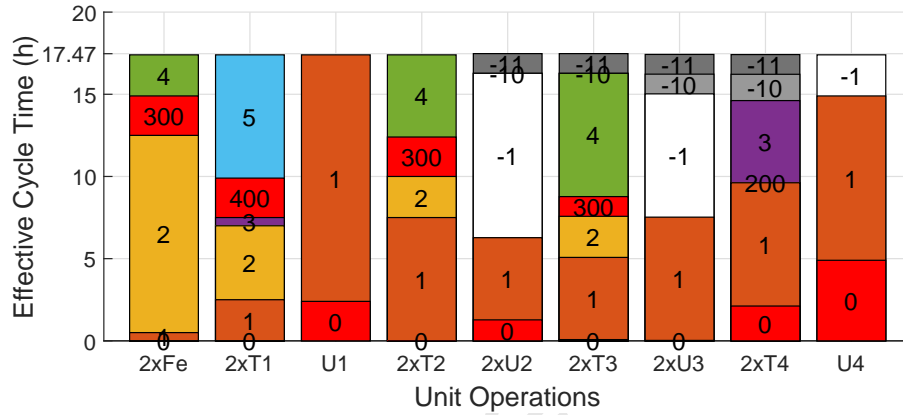
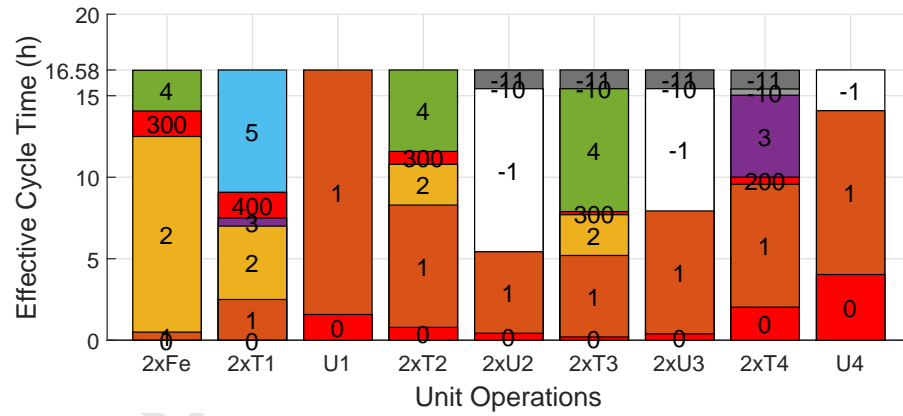


Figure (6) Excerpt of improved CIP schedule with reduced waiting time.



(a) Line capacity under standard CIP policy.



(b) Line capacity under improved CIP schedule.

Figure (7) Effective cycle times for a campaign of 300 batches. Step nomenclature: 0:idle, 1,2,3....:processing, 200,300,....:waiting, -1:reinitialisation, -10:blocked by CIP-call, -11:CIP in progress

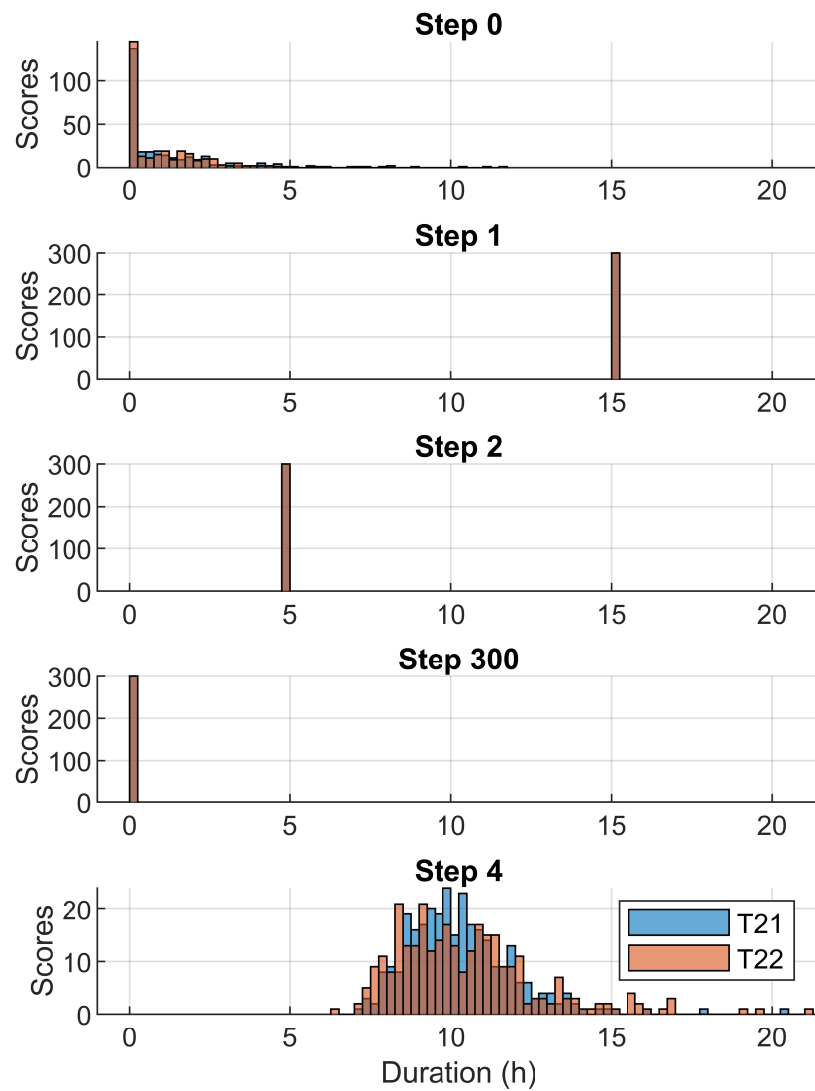


Figure (8) Durations of operations on tanks 21/22 as a consequence of the randomised flow rates on the downstream processing units.

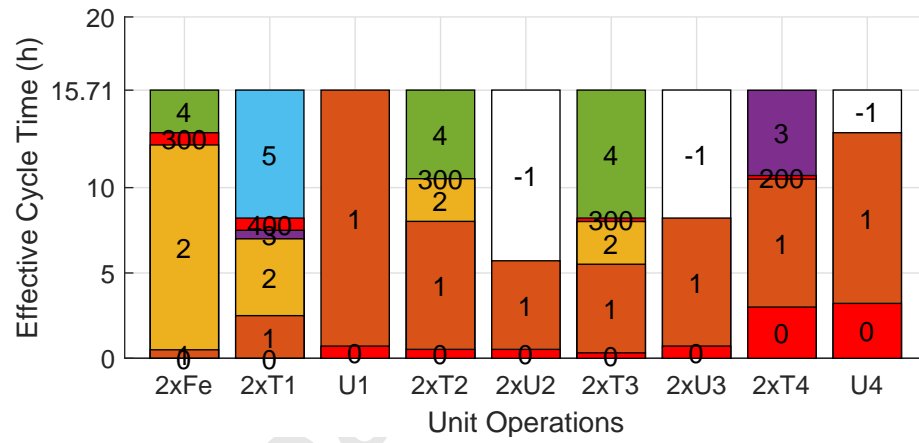


Figure (9) The equipment utilisation throughout the plant as a consequence of variability on the flow rates on U21/U22.